

M - matroid, ground set $\{n\}, E$

Definition

$$\text{trop}(M) := \left\{ w \in \mathbb{R}^n \mid \begin{array}{l} \forall C \in \mathcal{C}(M) \quad \exists i, j \in C, i \neq j \\ \text{s.t. } w_i = w_j = \min \{w_\ell \mid \ell \in C\} \end{array} \right\}$$

"minimum is attained twice"

Theorem

$$\boxed{\text{trop}(M) = \widetilde{\mathcal{S}}(M)}$$

Pf: First prove $w \in \text{trop}(M) \Leftrightarrow M_w$ loopless

Bases = bases of M
of max. w -weight

Note: $w \notin \text{trop}(M) \Leftrightarrow \exists C \in \mathcal{C}(M), \exists i \in C : w_i < w_j \quad \forall j \in C \setminus \{i\}$

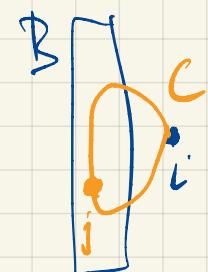
- If M_w has a loop, say i , then no basis of M_w contains i .

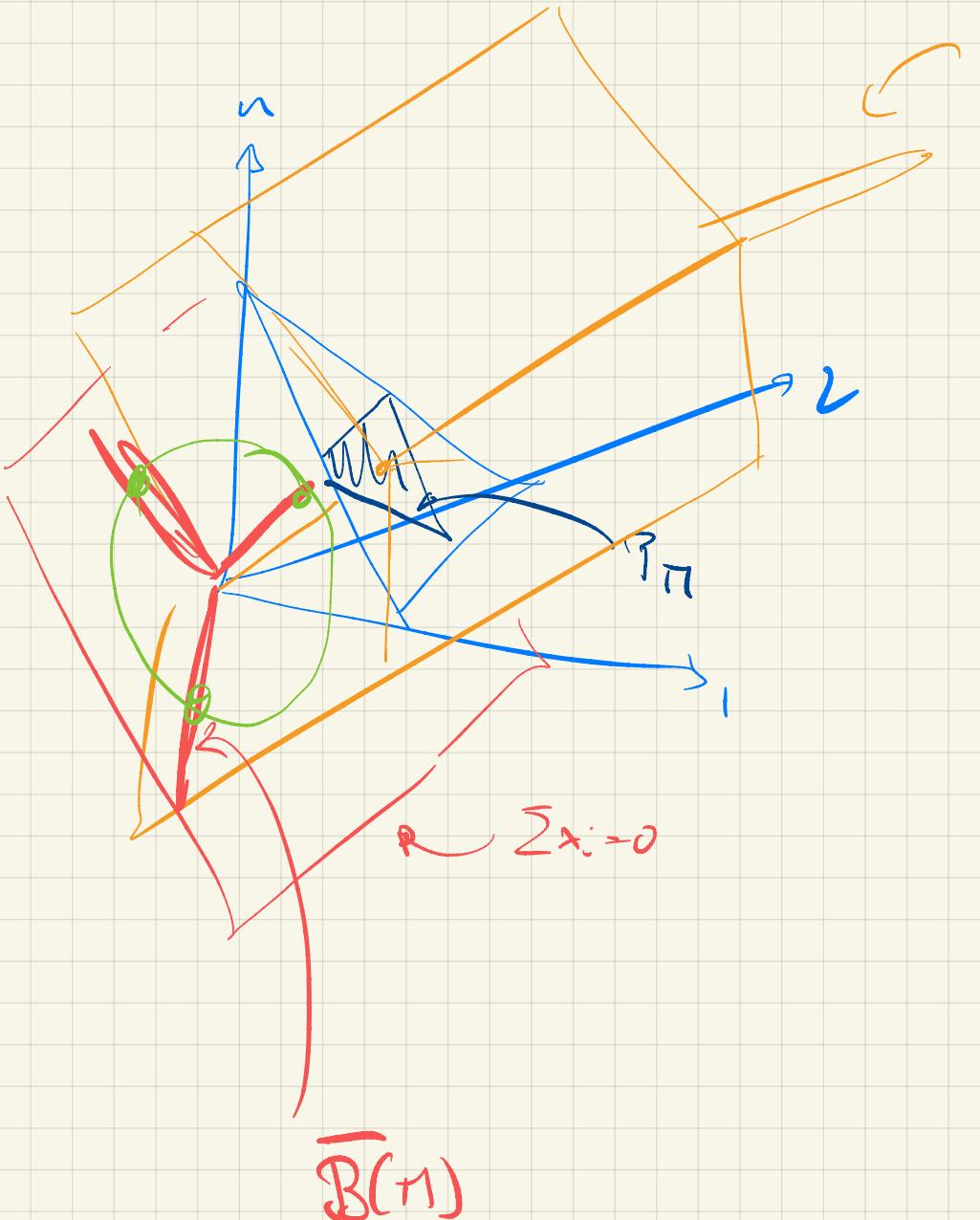
Take such a Basis. B . Then:

$$B \cup \{i\} \supseteq C \text{ for some circuit, } \underline{i \in C}$$

Claim: $w_i < w_j \quad \forall j \in C \setminus \{i\}$. In fact: if not, can exchange i for j and get basis of higher w -weight.

$B \cup \{j\} \cup \{i\}$ basis of M





$\tilde{B}(\pi) = w's \text{ with } \Gamma_w \text{ loops}$

$U_{3,5}$

" $B(\pi) \subseteq \overline{B}(\pi) \subseteq \tilde{B}(\pi)$ "

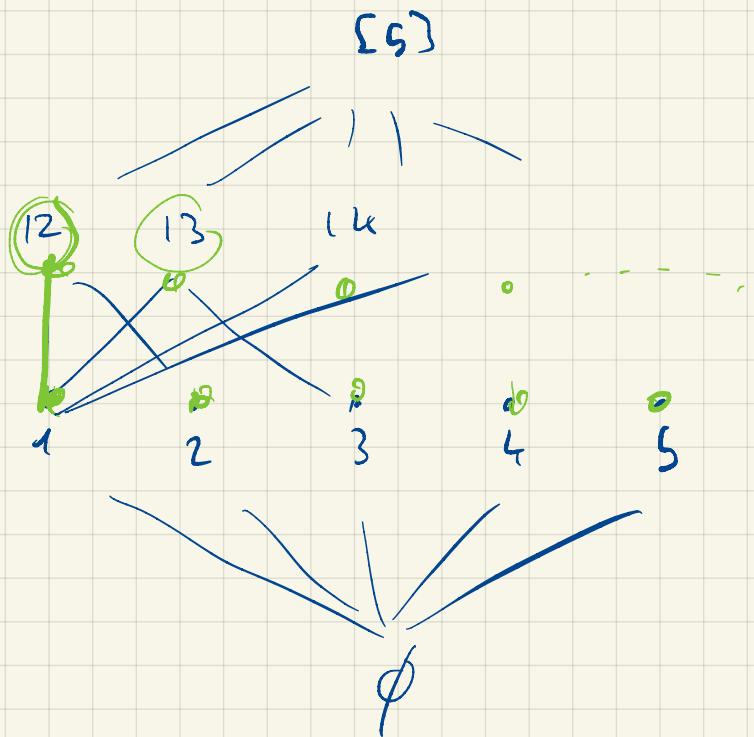


here: $B(\pi)$

For "computations":
flats

$B(\pi) = \overline{B}(\pi) \cap \text{unit-sphere}$

Example: $\pi = \underline{\cup}_{3,5}$

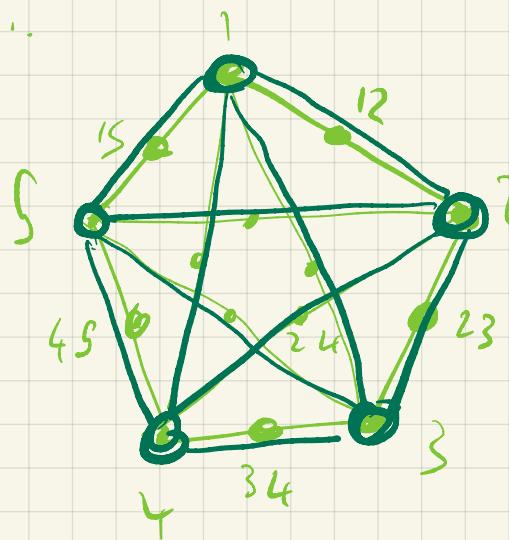


$\mathcal{B}(\pi)$:

Last line:

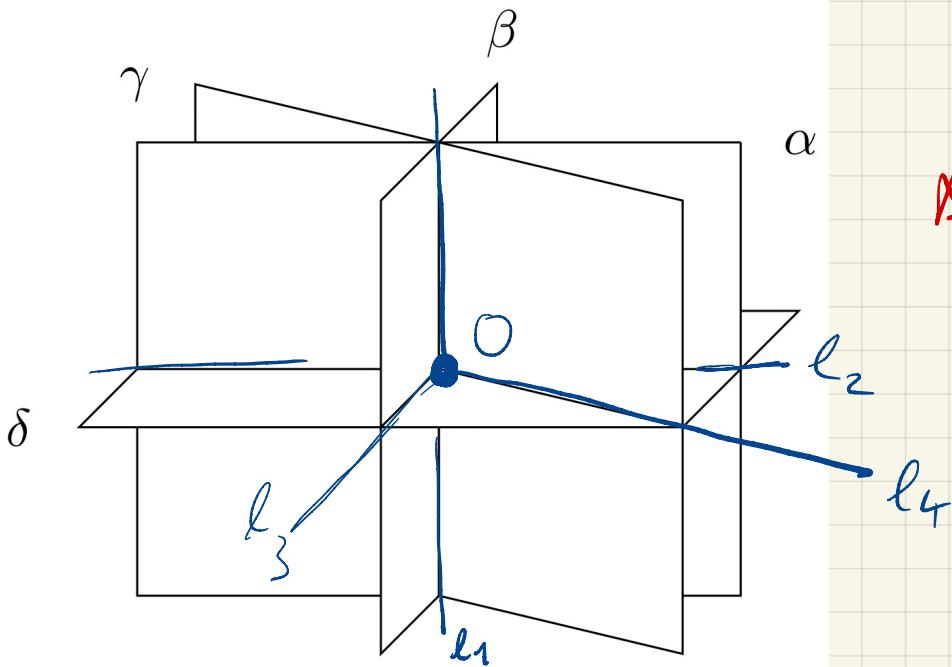
$\mathcal{B}(\pi)$ "is subdivided by" the couplet of chains of proper flats
 ↳ + minimal.
 ↳ + maximal.

Flats:



couplet of
chains of
flats.

A:



"Arrangement of hyperplanes"

Let V a K -vectorspace.

An. of $H.$ is a finite

$$A = \{H_1, \dots, H_n\}$$

{

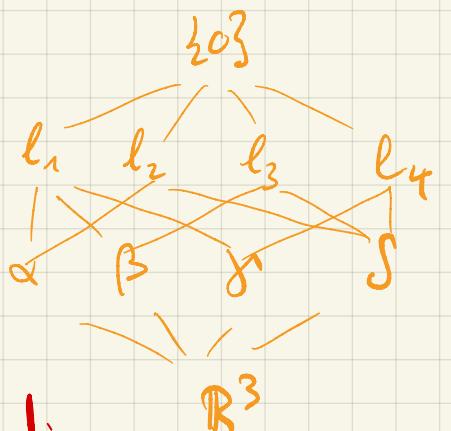
H_i is linear subspace
of codim. 1, i.e.

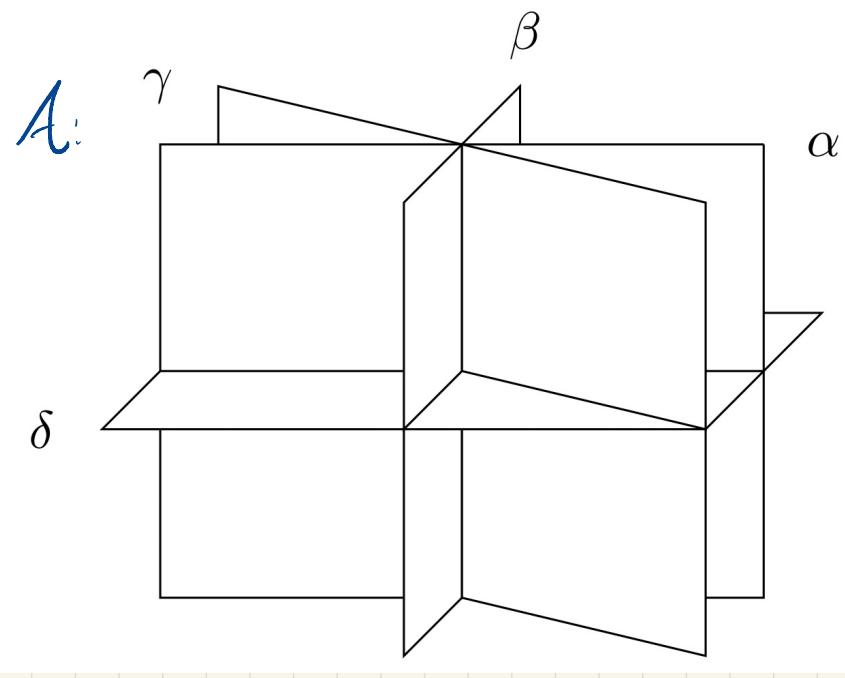
$$H_i = \{x \mid a_i^T x = 0\}$$

for some $a_i \in V$

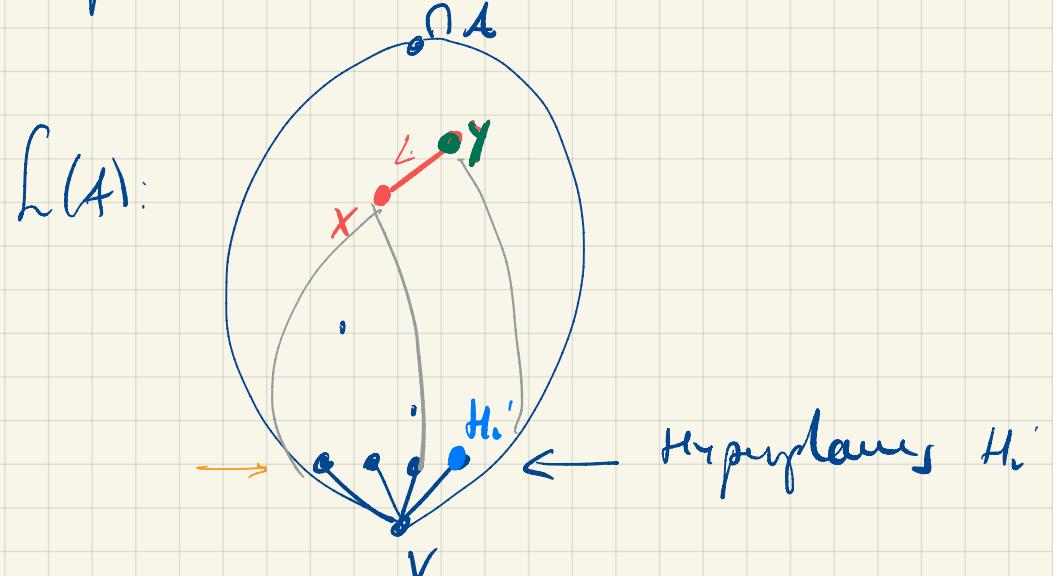
Poset of intersections:

$$\mathcal{L}(A) = \left\{ \bigcap_{i \in I} H_i \mid I \subseteq \{n\} \right\}, \quad X \leq Y \iff X \supseteq Y$$



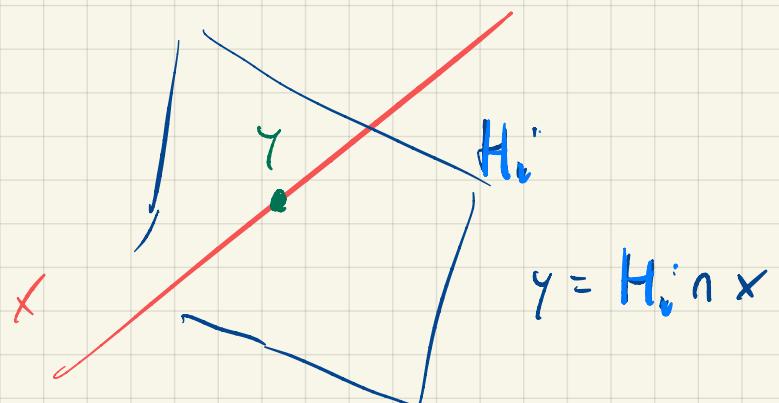


λ - ampt. of hyperplanes,
 $L(\lambda)$: poset of intersections



$H_i \not\leq x$ (otherwise $x \subseteq H_i$)

$H_i \leq y$



$x \leq y \iff \exists$ hyperpl. H_i , $H_i \not\leq x$, s.t. $y = H_i \cap x$

atom

" $x \leq z < y \Rightarrow x = z$ "

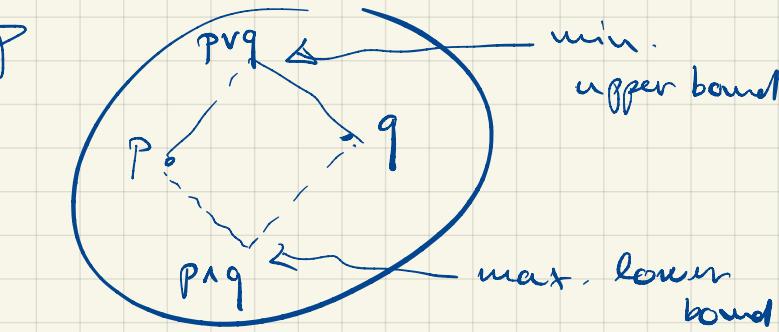
$H_i \nleq x$

"Partially ordered set" (poset) : (P, \leq)

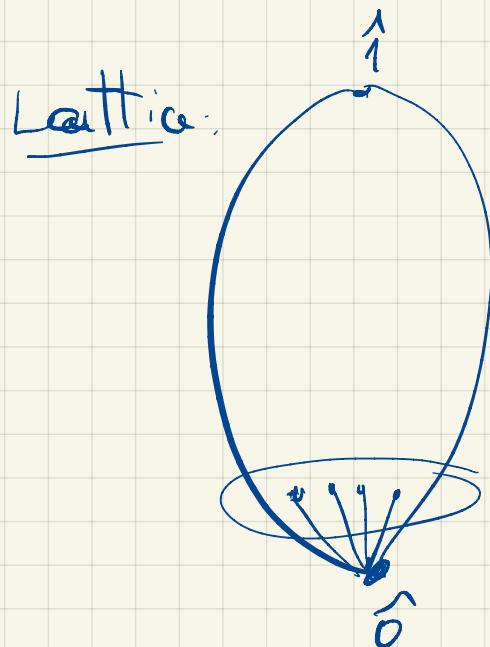
\nwarrow set ↗ relation
 (reflexive, antisymmetric, transitive)

Def: Call P a lattice if, $\forall p, q \in P$,

- $P_{\geq p} \cap P_{\geq q}$ has unique min. element, $p \vee q$, "join"
- $P_{\leq p} \cap P_{\leq q}$ " " max. " $p \wedge q$, "meet"



Note: Every finite lattice has a unique maximal ($\hat{1}$) & a unique minimal ($\hat{0}$) element



Example: $L(A)$ is a lattice with $x \vee y = x \wedge y$,
 $\hat{0} = V$, $\hat{1} = \cap A$.

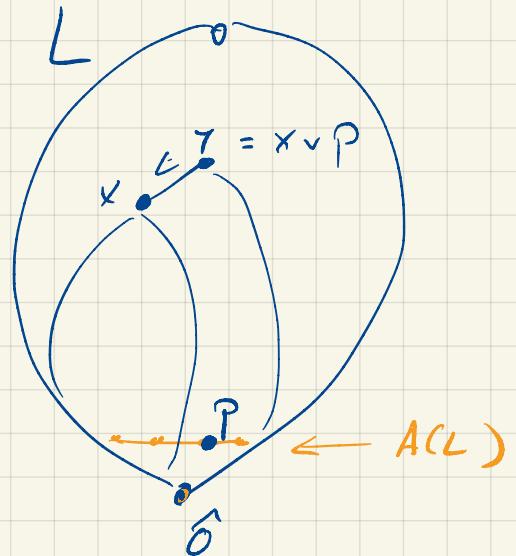
Def. If a poset P has $\hat{0}$, call atoms of P the elements of

$$A(P) := \{p \in P \mid p > \hat{0}\}$$

Definition Let L be a finite lattice.

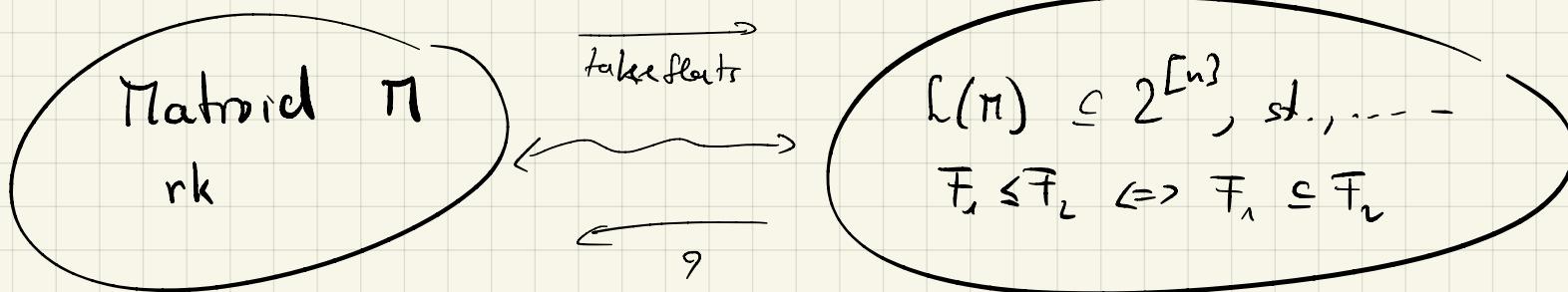
Call L a geometric lattice if, $\forall x, y \in L$

$$(G) \quad x < y \Leftrightarrow \exists p \in A(L), p \not\in x, \text{ s.t. } y = x \vee p$$



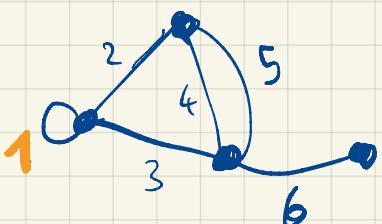
Example: $L(\pi)$.

→ Matroids "are" geometric lattices

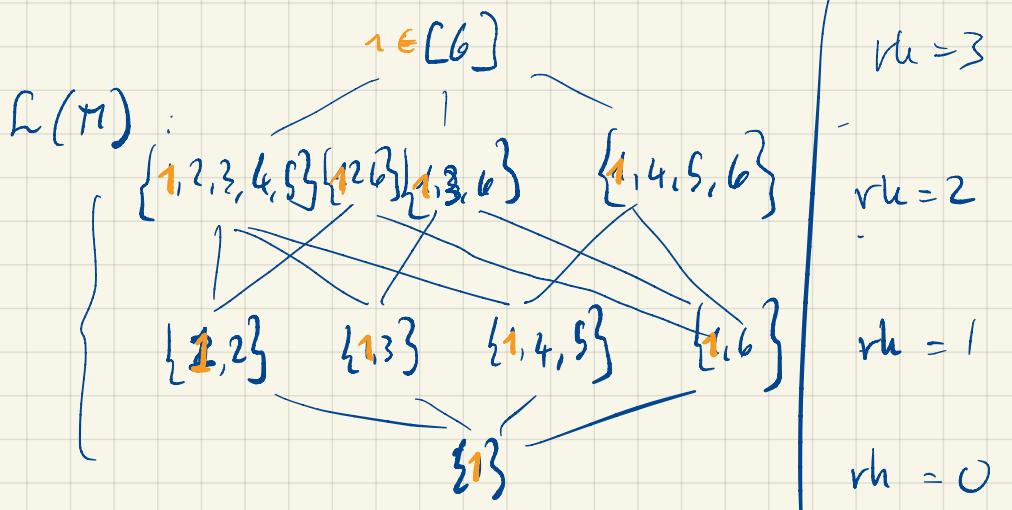


Example : Let $M = M(G)$

for G :



$$rh(\{3, 4\})$$



$L(M[2, 3, \dots, 6])$ poset-isom. to $L(\pi)$

$$rh(\{3, 4\}) = rh(\text{cl}(\{3, 4\})) = 2$$

Say M is simple (no loops & no parallel elements, like $\{4, 5\}$),

then atoms of $L(M)$ are single-element sets.

Then, given $X \subseteq [n]$ can read off $rh(X)$ by:

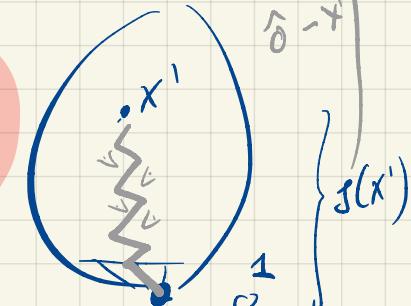
X

smallest $X' \in L(M)$
with $X \subseteq X'$

$$"X' = \text{cl}(X)"$$



look "on what
Floor" X' lies

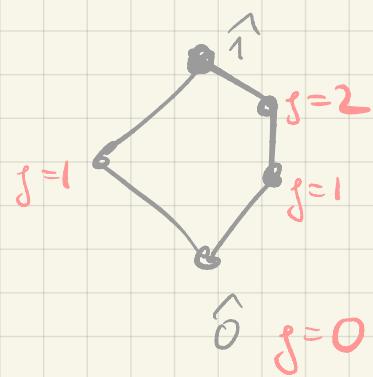


length of
any max.
chain
 $\hat{0} \dashv \hat{1}$

Def: Let P be poset. a rank function on P is

$g: P \rightarrow \mathbb{N}$ s.t. (i) $g(x) = 0$ if x is a minimal element

(ii) $g(x)+1 = g(y)$ if $x < y$



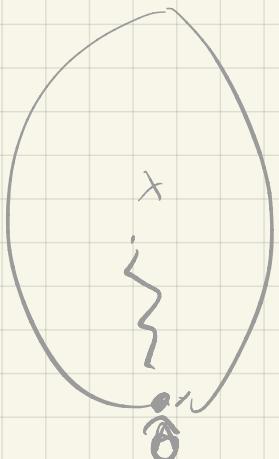
if a rank f. exists, call P ranked

(Notice: if P has $\hat{0}$ & a rank function g , then g is unique)

Def A chain in P is any subset $w = \{x_0 < x_1 < \dots < x_n\}$. Its length is $l(w) = |w| - 1$

Lemma 3.3 In a geometric lattice any two max. chains between the same elements have same length

(so: Geom. l. are ranked &c...)



$g(x)$ = length of any maximal $\hat{0} - x$ chain

Lemma.

Let L be a geometric lattice

with rank function $g \cdot \left(\begin{array}{l} g(\gamma) = \text{length} \\ \text{of max } \hat{0}-\gamma \text{ chain} \end{array} \right)$

Then, for all $x, y \in L$

$$g(x) + g(y) \geq g(x \wedge y) + g(x \vee y)$$

Proof: let $\stackrel{x \wedge y}{z_0} < z_1 < \dots < \stackrel{y}{z_k}$ saturated chain,

$$\text{then } k = g(y) - g(x \wedge y).$$

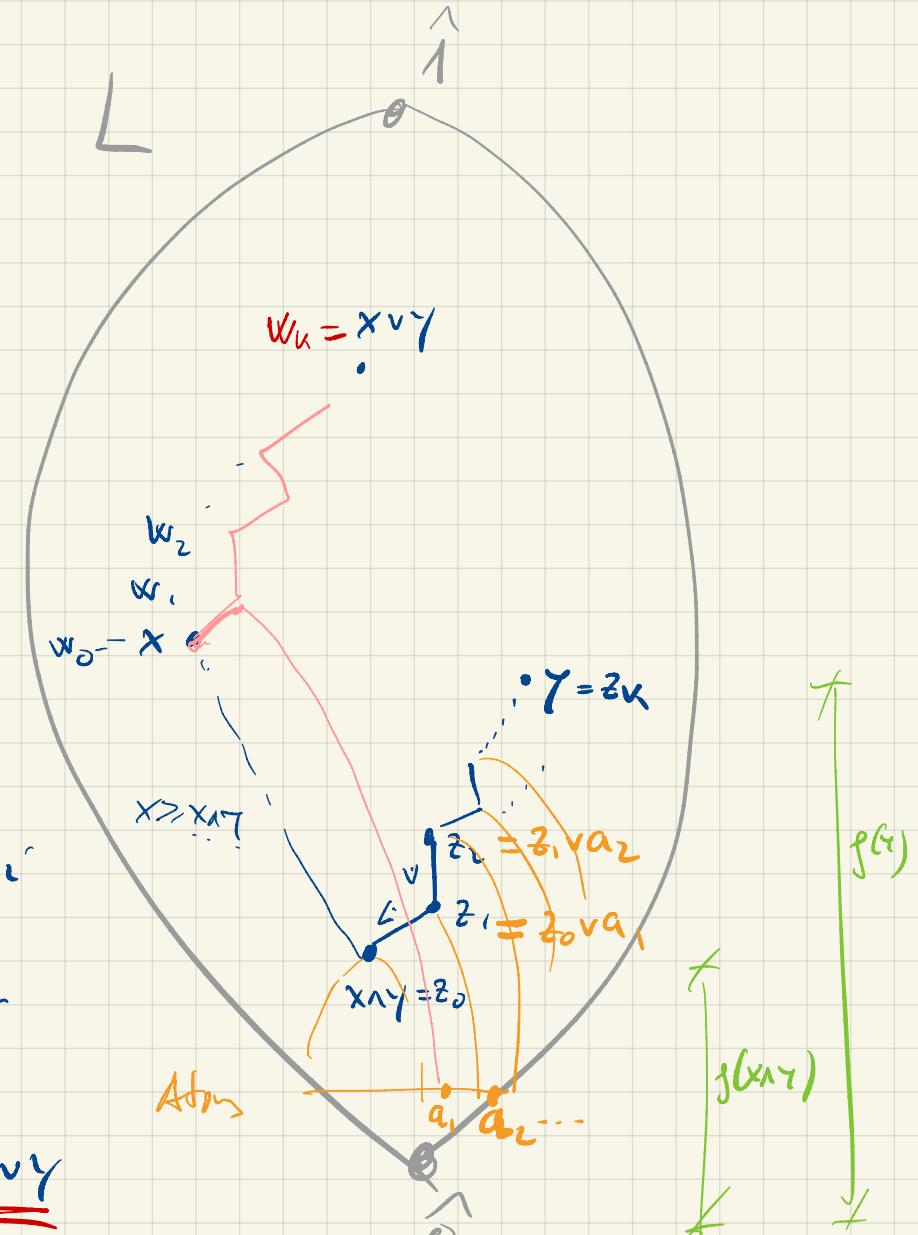
By (G) there are atoms a_1, a_2, \dots s.t. $z_i = z_{i-1} \vee a_i \quad \forall i$

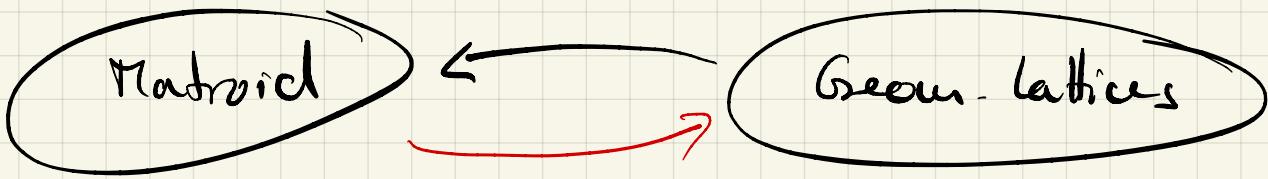
Define new $w_0 \leq w_1 \leq \dots \leq w_k$ by $w_0 = x$
 $w_i = w_{i-1} \vee a_i$

$$\stackrel{?}{=} w_k = x \vee a_1 \vee a_2 \dots \vee a_k = x \vee \underbrace{(x \wedge y)}_{z_k = y} \vee a_1 \vee \dots \vee a_k = \underline{\underline{x \vee y}}$$

By (G) for all i either $w_i = w_{i-1}$ or $w_i > w_{i-1}$

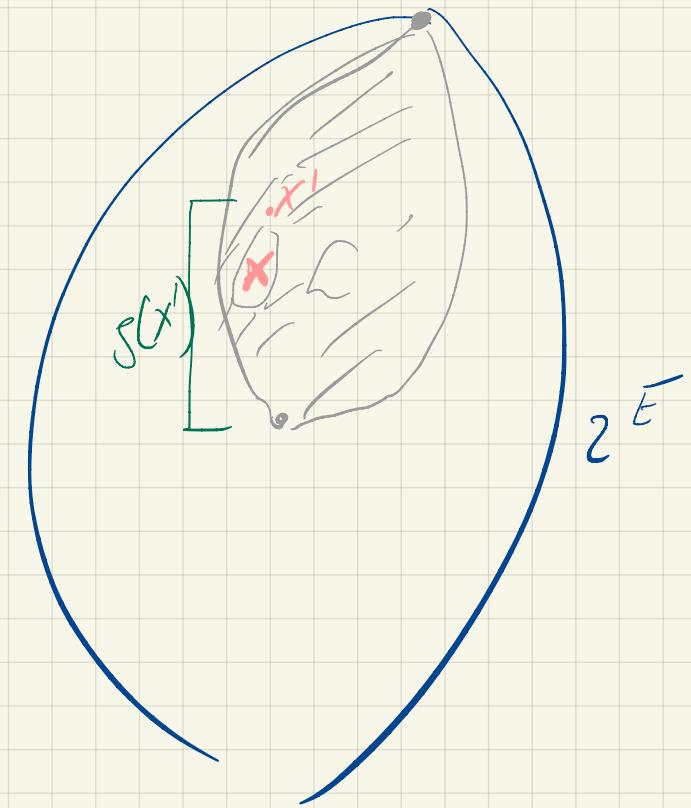
Thus $g(x \vee y) - g(x) = \text{length } \{w_0 \leq w_1 \leq \dots \leq w_k\} \leq k = \stackrel{\text{Ass.}}{=} g(y) - g(x \wedge y)$





Proposition 3.7. Let E be a finite set and let $\mathcal{L} \subseteq 2^E$ a family of subsets of E , partially ordered by inclusion and such that $E \in \mathcal{L}$. Suppose further that \mathcal{L} is a geometric lattice with rank function ρ , with meet operation given by set intersection, and such that the union of the atoms of \mathcal{L} equals E . Then, for every $X \subseteq E$ there is a unique minimal X' in \mathcal{L} such that $X \subseteq X'$, and the extension r of ρ on 2^E given by $r(X) := \rho(X')$ is a matroid rank function.

$$rk(x) := g(x')$$



For "Matroid \rightarrow Geom. lattice",

Given matroid M let $L(M)$ the set of flats, ordered by inclusion.

Then, for $F_1, F_2 \in L(M)$:

$$F_1 \vee F_2 = \text{cl}(F_1 \cup F_2)$$

$$F_1 \wedge F_2 = F_1 \cap F_2$$

← Lecture 2

So $L(M)$ is a lattice, and it is geometric (Prop. 3.13)

$A = \{H_1, \dots, H_n\}$ anyf. of hypers.,
say $H_i = \{a_i^T x = 0\}$

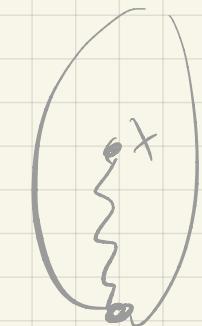
today



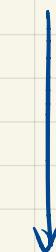
{(A) geom. lattice,
rank fn. g , induces
a matroid on $[n]$,
with rank function

r_{lat}

$$(r_{\text{lat}}(X) = \text{codim}(X))$$



$\{a_1, \dots, a_n\}$



long ago

Matroid of linear
dependencies, with
rank function

r_{dep}

$$(r_{\text{dep}}(X) = \dim \text{span}(X))$$

IN FACT:

$$r_{\text{lat}} \equiv r_{\text{dep}}$$

"the one and only matroid"

Ded Let \mathcal{A} be arrangement of hyperplanes in V . Then $\mathcal{M}(\mathcal{A}) := V \setminus \cup \mathcal{A}$
 (- interesting e.g. over \mathbb{C})

Remember. $\mathcal{A} = \{H_1 \dots H_n\}$, "normal vectors" $a_1 \dots a_n$, arrgt. in \mathbb{K}^d ,

$$A = \begin{bmatrix} | & | & | \\ a_1 & a_2 & \dots & a_n \\ | & | & \dots & | \end{bmatrix} = \begin{bmatrix} & b_1 & & \\ & \vdots & & \\ & b_d & & \end{bmatrix}$$

say $nA = \{0\}$

Then A is full-rank, and so is A^t

Now $A^t: \mathbb{K}^d \rightarrow \mathbb{K}^n$ is injective, with image
 A^t ↗ $V := \text{im } A^t = A^t \mathbb{K}^d$

Notice: $x \in H_i$ iff $a_i^t x = 0$, iff $(A^t x)_{\underline{i}} = 0 \Leftrightarrow \underline{\text{A}^t x \in V \cap \{y_{\underline{i}} = 0\}}$
 with coordinate plane

So: $x \in \mathcal{M}(\mathcal{A})$ iff $A^t x \in (V \setminus \cup \{y_{\underline{i}} = 0\}) = (\mathbb{K}^*)^n$

Can we characterize $V = \text{im } A^t$ as solution of set of equations?

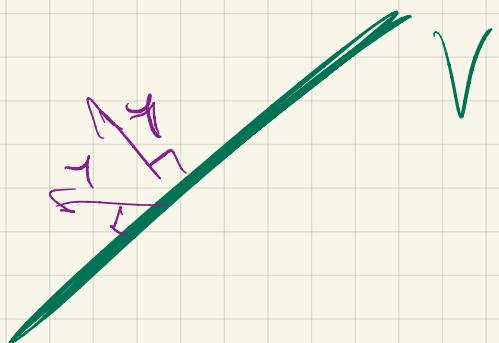
So look at $V = \text{im } A^t = \text{im} \begin{bmatrix} | & | \\ b_1 & b_2 \\ | & | \end{bmatrix}$

for $\gamma \in \mathbb{R}^n$

we have $\gamma^t V$ iff $\gamma \in \ker(A)$

$$A = \begin{bmatrix} | & | \\ a_1 & a_2 \\ | & | \end{bmatrix} \begin{bmatrix} \gamma_1 \\ \vdots \\ \gamma_n \end{bmatrix}, \text{ so } \gamma \in \ker(A) \text{ iff } \sum \gamma_i a_i = 0$$

in particular, if $\gamma \in \ker(A)$, then $\{\gamma_i \mid \gamma_i \neq 0\}$ is a dep. set
in the matroid $M(A)$



Ultimately:

$$V = \left\{ \sum \gamma_i x_i = 0 \mid \begin{array}{l} \text{for some } \gamma \in \ker(A) \\ (\text{a basis suffices}) \end{array} \right\}$$

a generating set for $\ker(A)$
is $\{\underline{v(C)} \mid C \text{ circuit of the } M(A)\}$

In fact if C is circuit, then there is a vector

$v(C) \in \ker(A)$ with $C = \{i \mid v(C)_i \neq 0\}$, and $v(C)$ unique up to scalar mult!

Proof:- let $\lambda_1 a_1 + \dots + \lambda_n a_n = 0$ minimal-support linear dependency.
 $\lambda_i \neq 0$
 (pertains to circuit $\{1, 2, \dots, n\}$)

Any other lin. dependency supported on $1, \dots, h$:

$$\mu_1 a_1 + \dots + \mu_h a_h = 0 \quad \mu_i \neq 0$$

Look at subtraction $\lambda_1 a_1 + \dots + \lambda_n a_n = 0$

$$\frac{\lambda_1}{\mu_1} \mu_1 a_1 + \frac{\lambda_1}{\mu_1} \mu_2 a_2 + \dots + \frac{\lambda_1}{\mu_n} \mu_n a_n = 0$$

$\Rightarrow \mu_1$ is a scalar multiple of λ_1 .

0

Must be all zero,
 otherwise this is
 lin. dependency
 "smaller" than C

Summary:

$$\mathcal{M}(A) \approx \bigvee_{\substack{\text{in } A^t}} \left(\mathbb{K}^* \right)^n$$

and \wedge is defined inside $(\mathbb{K}^*)^n$ by

$$\sum_i v(c)_i x_i = 0 \quad , \text{ one for each circuit } C$$

Now let $\mathbb{K} = \mathbb{C}$

$\underbrace{\sum_i}_{\substack{\text{taking} \\ \text{minim}}} \underbrace{v(c)_i}_{x_i} \underbrace{= 0}_{\substack{\text{"min. is attained finitely"}}$

$$\text{Trop}(\mathcal{M}(A)) = \underline{\text{trop}(\tau)} = \widetilde{\mathcal{S}}(\tau)$$

anything like this, where τ is any matroid,
is a tropical linear space.

Those that arise from $(\mathbb{C}-)$ representable
matroids are only some of them
(called "tropicalized linear spaces")

