# NOTES FOR LECTURE 2

This Lecture Notes follow the treatment of this subject given in the papers by Gel'fand-Goresky-MacPherson-Serganova and by Feichtner-Sturmfels that are includer in the literature folder.

# 1. MATROID POLYTOPES

1.1. **Basic definitions.** Let  $n \in \mathbb{N}$ . The *standard-simplex*  $\Delta$  in  $\mathbb{R}^n$  is the convex hull of the standard basis vectors,

$$\Delta := \operatorname{conv} \{ e_i \mid i = 1, \dots, n \} \subseteq \mathbb{R}^n.$$

We write  $\Delta^{(n)}$  if specification is needed. Notice that every edge of  $\Delta$  is parallel to some vector of the form  $e_i - e_j$  for  $i, j \in [n]$ .

Recall that we write  $[n] := \{1, ..., n\}$  and, for  $r, n \in \mathbb{N}$ ,  $\binom{n}{r} := \{A \subseteq [n] \mid |A| = r\}$ .

Moreover, for any subset  $A \subseteq [n]$  we write

$$e_A := \sum_{i \in A} e_i \in \mathbb{R}^n$$

for the sum of the basis vectors indexed by elements of A.

**Definition 1.1.** Given any family  $\mathcal{X} \subseteq 2^{[n]}$ , we consider then the polytope

$$P_{\mathcal{X}} := \operatorname{conv}\{e_X \mid X \in \mathcal{X}\}.$$

# 1.2. The main theorem.

**Theorem 1.2** (Gel'fand-Goresky-Macpherson-Serganova). A nonempty family  $\mathcal{B} \subseteq {n \choose r}$  is the set of bases of a rank r matroid on the ground set [n] if and only if all edges of  $P_{\mathcal{B}}$  are parallel to edges of  $\Delta$ .

**Remark 1.3.** Often the statement of Theorem 1.2 is formulated in terms of a subset of vertices of the "hypersimplex"  $P_{\binom{n}{r}}$  rather than "a nonempty family  $\mathcal{B} \subseteq \binom{n}{r}$ ". The equivalence is obvious.

We comment on one aspect of the alternative formulation that is not immediatly apparent in ours: if  $\mathcal{B} \subseteq \binom{n}{r}$ , then the set of vertices of  $P_{\mathcal{B}}$  is exactly  $\{e_B \mid B \in \mathcal{B}\}$ . In order to see this, given  $X \in \binom{n}{r}$  consider the linear functional  $\ell_X : \mathbb{R}^n \to \mathbb{R}, x \mapsto \langle e_X \mid x \rangle$ . Then,  $\ell_X(e_Y) \leq |X|$  for all  $Y \in \binom{n}{r}$ , with equality if and only if X = Y. In particular,  $e_X$  cannot be a convex combination of any subset of the  $e_Y, Y \in \binom{n}{r} \setminus \{X\}$ .

**Definition 1.4.** Given a matroid M we will write  $P_M$  for  $P_{\mathcal{B}(M)}$ . Every such polytope is called a *matroid polytope*.

*Proof of Theorem 1.2.* We prove only one direction, and leave the other as a reading assignment, from the original paper (§4.3 of the paper by Gel'fand, Goresky, MacPherson and Serganova, available in the literature folder).

Assume then that every edge of the nonempty polytope  $P_{\mathcal{B}}$  is parallel to a vector of the form  $e_i - e_j$  for some  $i, j \in [n]$ . We have to prove that  $\mathcal{B}$  satisfies axiom ( $\mathcal{B}2$ ). Let then  $B_1, B_2 \in \mathcal{B}$ . Up

 $P_M$  (Matroid polytope)

to re-ordering the coordinates we can assume that

$$e_{B_1} = (1, \dots, 1, 0, \dots, 0, 1, \dots, 1, 0, \dots, 0)$$
  

$$e_{B_2} = (\underbrace{0, \dots, 0}_{A\text{-coord.}}, \underbrace{1, \dots, 1}_{B\text{-coord.}}, \underbrace{1, \dots, 1}_{C\text{-coord.}}, \underbrace{0, \dots, 0}_{D\text{-coord.}})$$

For some partition  $[n] = A \uplus B \uplus C \uplus D$ .

Now, by convexity of  $P_{\mathcal{B}}$  the segment between  $e_{B_1}$  and  $e_{B_2}$  is fully contained in  $P_{\mathcal{B}}$ , which is contained in the vertex cone  $C(P_{\mathcal{B}}, v)$  of  $P_{\mathcal{B}}$  at  $v := e_{B_1}$ . The vertex cone was treated in the Warm-up to this lecture, where it was also explained that this cone is generated by the directions of the edges emanating from v, and is based at v. In particular,

$$e_{B_2} - e_{B_1} = (-1, \dots, -1; 1, \dots, 1; 0, \dots, 0; 0, \dots, 0) = \sum_{\substack{\text{conv}\{v, w\} \text{ is} \\ \text{edge of } P_{\mathcal{B}}}} \lambda_w(w - v) \tag{\dagger}$$

with all  $\lambda_w \in \mathbb{R}_{\geq 0}$ . Now by assumption for every vertex w of  $P_{\mathcal{B}}$  appearing in (†) the vector w - vis a positive multiple of  $e_i - e_j$  for some  $i, j \in [n]$ , i.e., in coordinates i and j we have  $w_i > v_i$ and  $w_j < v_j^{1}$ . Now both v and w are 0 - 1 vectors (as is the case for every vertex of  $P_{\mathcal{B}}$ ), thus necessarily  $v_i = (e_{B_1})_i = 0$  and  $v_j = (e_{B_1})_j = 1$ . This means that  $i \in B \cup D$  and  $j \in A \cup C$ .

In particular, on the r.h.s of (†) there is never a negative contribution to coordinates in D. This implies that, if  $\lambda_w > 0$ , then  $i \notin D$ , since otherwise we'd have a positive contribution of  $e_i$  to a D-coordinate that can't be cancelled by other terms. Analogously,  $j \notin C$  whenever  $\lambda_w > 0$ .

In summary, for every vertex w adjacent to v such that  $\lambda_w > 0$  in  $(\dagger)$ , if  $e_i - e_j$  is the direction of the edge between w and v then  $i \in B$  and  $j \in A$ .

Now we return to checking ( $\mathcal{B}2$ ). Let  $b \in B_1 \setminus B_2$  and recall that in fact  $B_1 \setminus B_2 = A$ . Then  $(e_{B_1} - e_{B_2})_b = -1$ , and so there must be w with  $\lambda_w > 0$  in ( $\dagger$ ) such that w - v is a positive multiple of  $e_{b'} - e_b$  for some  $b' \in B = B_2 \setminus B_1$ . Since w and v are vertices of  $P_{\mathcal{B}} \subseteq \{x \mid 0 \leq |x_k| \leq 1 \text{ for all } k\}$ , in fact  $w - v = e_{b'} - e_b$ . Thus, for the  $B' \in \mathcal{B}$  such that  $w = e_{B'}$  we have

$$B' = (B_1 \setminus \{b\}) \cup \{b'\},$$

and the proof is complete.

Some corollaries follow directly from Theorem 1.2.

**Corollary 1.5.** Every face of a matroid polytope is a matroid polytope.

**Corollary 1.6.** Every matroid polytope is a subset of the simplex  $r \cdot \Delta$ , where r is the rank of the matroid and  $\Delta = \Delta^{(n)}$ , with n the cardinality of the ground set. Notice that  $r \cdot \Delta$  is described by  $x_1 \ge 0, \ldots, x_n \ge 0$  and  $\sum_i x_i = r$ .

1.3. Flats and flacets. Our next goal is to determine which inequalities determine a matroid polytope as a subset of  $r \cdot \Delta$ . To this end we need a definition.

**Definition 1.7.** Let M be a matroid with rank function  $\mathrm{rk}$  on the ground set E. A subset  $F \subseteq E$  is called a *flat* of M if  $\mathrm{rk}(F \cup \{e\}) > \mathrm{rk}(F)$  for all  $e \in E \setminus F$ .

**Lemma 1.8.** Let  $F_1, F_2$  be flats of M. Then,  $F_1 \cap F_2$  is a flat of M.

*Proof.* Let  $e \in E \setminus (F_1 \cap F_2)$ . In particular, there is  $j \in \{1, 2\}$  for which  $e \notin F_j$ . Then,

 $\operatorname{rk}(F_j) + \operatorname{rk}((F_1 \cap F_2) \cup \{e\}) \ge \operatorname{rk}(F_1 \cap F_2) + \operatorname{rk}(F_j \cup \{e\}).$ 

Since  $F_j$  is a flat,  $\operatorname{rk}(F_j) < \operatorname{rk}(F_j \cup \{e\})$ , and the inequality above implies  $\operatorname{rk}(F_1 \cap F_2) \cup \{e\}) > \operatorname{rk}(F_1 \cap F_2)$  as desired.

(Flats)

<sup>&</sup>lt;sup>1</sup>I owe this formulation to Tim - thanks!



FIGURE 1. Running example for this lecture

**Corollary 1.9.** For every  $A \subseteq E$  there is a unique minimal flat of M, called cl(A) ("closure" of A), that is inclusion-minimal among all flats containing A. The closure satisfies rk(A) = rk(cl(A)).

Proof. Set

(Closure)

$$\mathrm{cl}(A):=\bigcap_{\substack{F \text{ flat of } M\\ A\subseteq F}}F$$

and apply Lemma 1.8 to obtain the first claim.

For the second claim, consider  $\overline{A} := \max\{X \supseteq A \mid \operatorname{rk}(X) = \operatorname{rk}(A)\}$ . This definition is wellposed: indeed, given X', X'' in the set on the r.h.s., (R3) applied to X' and X'' implies that  $\operatorname{rk}(X' \cup X'') = \operatorname{rk}(A)$ , hence  $X' \cup X''$  is of the set on the r.-h. s. as well, and so the maximum exists. Now, clearly  $\overline{A}$  is a flat and  $\operatorname{rk}(\overline{A}) = \operatorname{rk}(A)$ . Thus  $\operatorname{cl}(A) \subseteq \overline{A}$ , which implies the middle inequality of the following expression

$$\operatorname{rk}(A) = \operatorname{rk}(\overline{A}) \ge \operatorname{rk}(\operatorname{cl}(A)) \ge \operatorname{rk}(A)$$

where the last inequality holds by monotonicity of the rank function.

Now obviously all above inequalities must be equalities, and the second part of the claim follows.  $\Box$ 

**Theorem 1.10.** Let M be a matroid on the ground set [n] with rank function rk. Then

$$P_M = \{ x \in r \cdot \Delta^{(n)} \mid \sum_F x_i \le \operatorname{rk}(F) \text{ for all flats } F \text{ of } M \}$$

Proof.

Claim 1 For every  $A \subseteq E$ ,  $P_M$  satisfies the inequality

$$\sum_{i \in A} x_i \le \operatorname{rk}(A). \tag{1}$$

Moreover, the inequality is satisfied with equality for at least one vertex of  $P_M$ . *Proof.* We prove that every vertex of  $P_M$  satisfies (1). Let B be a basis of M and consider the associated vertex of  $P_M$  given by  $v := e_B$ . Then,  $\sum_{i \in A} v_i = |B \cap A|$ , and  $B \cap A$  is independent. The maximum size of an independent subset of A is, by definition,  $\operatorname{rk}(A)$ , whence v satisfies (1). If we pick any maximal independent set  $I \subseteq A$  and extend I to a basis B of  $M^2$ , we have  $\sum_{i \in A} (e_B)_i = |B \cap A| = |I| = \operatorname{rk}(A)$ , thus the vertex  $e_B$  satisfies Equation 1 with equality.

Claim 2 Every facet Q of  $P_M$  that intersects the interior of  $r \cdot \Delta$  can be defined by an inequality of the form

$$\sum_{i \in A} x_i \le \operatorname{rk}(A)$$

for some  $A \subseteq E$ .

*Proof.* The claim is void, hence trivially true, in case dim  $P_M \leq 1$ . In that case the facets of  $P_M$  are its vertices, and those all lie on the boundary of the simplex  $r \cdot \Delta$ . Assume then dim  $P_M \geq 2$ .

First, we prove that we can find a normal vector to Q with all entries either equal to 0 or 1. A vector is normal to Q if and only if it is normal to the direction of every edge of Q. Now, the edges of Q are edges of  $P_M$  and so they have the direction  $e_i - e_j$  for some  $i, j \in [n]$ . Consider now on [n] the equivalence relation defined as the transitive closure of  $i \sim_Q j$  if Q has an edge in direction  $\pm (e_i - e_j)$ . Let X be an equivalence class of  $\sim_Q$  with |X| > 1. Such an A exists because Q has at least one edge (since dim  $Q = \dim P_M - 1 \geq 1$ ). Pick a vertex v of Q. Since dim  $P_M > \dim Q$  and  $P_M \subseteq C(P_m, v)$ , there must be an edge of  $P_M$  that does not lie in aff(Q). Letting  $e_k - e_l$  be the direction of this edge, we must have  $k \not\sim_Q l$  (since  $i \sim_Q j$  implies that aff(Q) contains an affine line in direction  $e_i - e_j$ ). Thus, up to relabeling k and l we can assume that the equivalence class  $[l]_{\sim_Q}$  is disjoint from X. Let then

$$A := X \cup [k]_{\sim_Q}, \qquad a := e_A.$$

This vector is orthogonal to every edge of Q (since  $a_i = a_j$  whenever  $i \sim_Q j$ ), but not to every edge of  $P_M$  (because  $a_k \neq a_l$ ). This means that the linear form  $\sum_{i \in A} x_i$  is constant on Q-say with value b-but not on  $P_M$ . In particular, since Q is a facet of the polytope  $P_M$  we know that the facet Q is defined either by  $\sum_{i \in A} x_i \geq b$  or  $\sum_{i \in A} x_i \leq b$ . If the former case arises, we can use the fact that  $P_M$  lies in the affine hyperplane  $\sum_{i \in E} x_i = r$  and substract r on both sides of the inequality, obtaining  $-\sum_{i \in E \setminus A} x_i \geq b - r$ , i.e.,  $\sum_{i \in E \setminus A} x_i \leq r - b$ . In any case (i.e., up to switching A with  $E \setminus A$  and b with r - b) there is a subset  $A \subseteq E$ 

In any case (i.e., up to switching A with  $E \setminus A$  and b with r-b) there is a subset  $A \subseteq E$ and a number  $b \in \mathbb{R}$  such that  $\sum_{i \in A} x_i \leq b$  is an inequality defining the facet Q. Now by Claim 1 we know that it must be  $b = \operatorname{rk}(A)$ .

<sup>&</sup>lt;sup>2</sup>This is possible by repeated application of  $(\mathcal{I}3)$ .

**Illustrative example.** Consider the example in Figure 1 and let Q be the "top horizontal" triangular facet (i.e., with vertices 13, 14, 12). If we want to determine a defining inequality for Q following the method of the proof of Claim 2, we first look at the equivalence relation  $\sim_Q$  which, in this case, has equivalence classes  $\{1\}, \{2, 3, 4\}$ . We would choose  $X = \{2, 3, 4\}$ . We'd then look at the direction of an edge exiting Q, say the edge between 12 and 24 with direction  $e_4 - e_1$ , hence k = 4, l = 1 so that  $[l]_{\sim_Q} = \{1\}$  is disjoint from X. Now, in this case  $[k]_{\sim_Q} = [4]_{\sim_Q} = X$ , therefore  $a = e_X = e_A = (0, 1, 1, 1)$ . Now  $\sum_{i \in A} x_i = x_2 + x_3 + x_4$  has value 1 on every vertex of Q, and value 2 on other vertices of  $P_M$ , thus we obtain the facet-defining inequality is  $x_2 + x_3 + x_4 \ge 1$ . This is not yet of the desired form. We substract  $x_1 + x_2 + x_3 + x_4 = 2 = r$  on both sides and we obtain  $-x_1 \ge -1$ , i.e.,  $x_1 \le 1$ , which is now of the desired form.

Claim 3 For every  $A \subseteq E$  and every  $x \in r \cdot \Delta$ , the inequality  $\sum_{i \in cl(A)} x_i \leq rk(cl(A))$  implies  $\sum_{i \in A} x_i \leq rk(A)$ .

 $\overline{Proof.}$  In fact,

$$\sum_{i \in A} x_i \le \sum_{i \in \operatorname{cl}(A)} x_i \le \operatorname{rk}(\operatorname{cl}(A)) = \operatorname{rk}(A).$$

The first inequality because  $A \subseteq cl(A)$ , the last equality by the second claim in Corollary 1.9.

Now, by Claims 1 and 2 the Theorem's claimed equality holds without restriction on F, and Claim 3 shows that it is possible to restrict F to be a flat.

#### 2. Minors, matroid connectivity and faces of $P_M$

2.1. Connectivity and dimension. Let M be a matroid on the ground set E. Define a relation on E via

 $e \sim_M f \Leftrightarrow$  either e = f, or there are bases  $B_1, B_2 \in \mathcal{B}(M)$  with  $B_2 = (B_1 \setminus \{e\}) \cup \{f\}$ .

**Lemma 2.1.** Let M be a matroid on the ground set E. Then, for all  $e, f \in E$ :

 $e \sim_M f \Leftrightarrow$  either e = f, or  $\{e, f\} \subseteq C$  for some circuit  $C \in \mathcal{C}(M)$ .

In particular,  $\sim_M$  is an equivalence relation.

*Proof.* Consider two elements  $e, f \in E, e \neq f$ .

If  $\{e, f\} \subseteq C$  for some circuit C, the set  $C \setminus \{f\}$  is independent and can be completed to a basis  $B_1$  of M not containing f. Now let  $B_2 := B_1 \setminus \{e\} \cup \{f\}$ . If there is any circuit  $C' \subseteq B_2$ , then it must contain f (otherwise  $C' \subseteq B_1$ , contradicting independence of  $B_1$ ). Now (C3) applied to C, C' and f would give a circuit fully contained in  $(C \cup C') \setminus \{f\} \subseteq B_1$ , a contradiction. Thus,  $B_2$  is independent and, since  $|B_1| = |B_2|$ , it is a basis.

On the other hand, if there are bases  $B_1, B_2 \in \mathcal{B}(M)$  with  $B_2 = B_1 \setminus \{e\} \cup \{f\}$ , then in particular  $B_1 \cup \{f\}$  is dependent, and there is a circuit  $C \subseteq B_1 \cup \{f\}$  with  $f \in C$ . If  $e \notin C$ , then  $C \subseteq B_2$ , a contradiction. Therefore  $\{e, f\} \subseteq C$  as desired.

That  $\sim_M$  is an equivalence relation follows from its formulation in terms of circuits by (C3).  $\Box$ 

c(M) (connectivity)

**Definition 2.2.** The equivalence classes of  $\sim_M$  are called *connected components* of M. The *connectivity* of the matroid M is the number c(M) of equivalence classes of  $\sim_M$ . We call M *connected* if c(M) = 1, disconnected otherwise.

# Remark 2.3.

(1) Since equivalence classes are always nonempty, for the empty matroid  $M = (\emptyset, \{\emptyset\})$  we have c(M) = 0, thus the empty matroid is disconnected according to Definition 2.2.

(2) It is important to notice that the notion of connectivity for matroids is *not* a generalization of the notion of connectivity for graphs (rather, if G is a graph, connectivity of M(G) has to do with "2-connectedness" of G, see §4.1 in Oxley's book). Can you find a connected graph G such that M(G) is disconnected as a matroid?

The notion of connectivity is directly related to the dimension of the matroid polytope.

**Proposition 2.4.** Let M be a matroid on the ground set [n]. Then

$$\dim P_M = n - c(M)$$

*Proof.* Lemma 2.1 shows that, for any  $i \neq j$ ,  $i \simeq_M j$  if and only if there are vertices  $e_{B_1}$  and  $e_{B_2}$  of  $P_M$  that are connected by a segment in the direction of  $e_i - e_j$ . In particular, the linear space in  $\mathbb{R}^n$  that is parallel to the affine span of  $P_M$  is generated by

$$S := \{e_i - e_j \mid i \sim_M j\} = \biguplus_{K \in E/\sim_M} \underbrace{\{e_i - e_j \mid i, j \in K\}}_{=:S_K}$$

where K runs over all connected components of M. Now the dimension of  $P_M$  is the (column-)rank of the matrix N whose columns are the elements of S. If we number  $K_1, \ldots, K_{c(M)}$  the connected components of M and rearrange the coordinate labels so that elements of  $K_h$  come before elements of  $K_k$  whenever h < k, the matrix N has the block-diagonal form

$$N = \begin{pmatrix} N_{1} & 0 & 0 & 0 \\ \hline 0 & N_{2} & 0 & 0 \\ \hline 0 & 0 & \ddots & 0 \\ \hline 0 & 0 & 0 & N_{c(\pi)} \\ \hline S_{K_{1}} & \cdots & S_{K_{c(n)}} \end{pmatrix}$$

Where, up to permuting and negating columns, the block  $N_k$  can be written

$$N_{k} = \begin{cases} 1 & 1 & 1 & 1 \\ -1 & 0 & 0 & 0 \\ 0 & -1 & 0 &$$

and has thus rank  $|K_k| - 1$ . It follows that the rank of N, and hence the dimension of  $P_M$ , equals  $\sum_k (|K_k| - 1) = n - c(M)$ .

# 2.2. Direct sums and cartesian products.

**Lemma 2.5.** Let  $M_1 = (E_1, \mathcal{I}_1)$ ,  $M_2 = (E_2, \mathcal{I}_2)$  be two matroids given in terms of independent sets and with disjoint ground sets  $(E_1 \cap E_2 = \emptyset)$ . Then

$$\mathcal{I}_1 \oplus \mathcal{I}_2 := \{ I_1 \cup I_2 \mid I_1 \in \mathcal{I}_1, I_2 \in \mathcal{I}_2 \}$$

is the family of independent sets of a matroid on  $E_1 \uplus E_2$ .

*Proof.* Exercise (see Worksheet 2).

Direct sum

**Definition 2.6.** Let  $M_1 = (E_1, \mathcal{I}_1)$ ,  $M_2 = (E_2, \mathcal{I}_2)$  be as in Lemma 2.5. The matroid  $M_1 \oplus M_2 := (E_1 \oplus E_2, \mathcal{I}_1 \oplus \mathcal{I}_2)$  is called the *direct sum* of  $M_1$  and  $M_2$ .

**Remark 2.7.** A matroid is disconnected (in the sense of  $\S2.1$ ) if and only if it can be written as a direct sum of two nonempty matroids. See Worksheet 2 for a precise statement.

**Corollary 2.8.** Let  $M_1$ ,  $M_2$  be matroids on disjoint ground sets. Then

$$\mathcal{B}(M_1 \oplus M_2) = \{B_1 \uplus B_2 \mid B_1 \in \mathcal{B}(M_1), B_2 \in \mathcal{B}(M_2)\}.$$

**Proposition 2.9.** Let  $M_1$ ,  $M_2$  be matroids on disjoint ground sets. Then

$$P_{M_1 \oplus M_2} = P_{M_1} \times P_{M_2}$$

Proof. From our Warm-up we know that  $P_{M_1} \times P_{M_2} = \operatorname{conv} \{ e_{B_1} + e_{B_2} \mid B_1 \in \mathcal{B}(M_1), B_2 \in \mathcal{B}(M_2) \}$ . Now the claim follows from Corollary 2.8 since  $e_{B_1} + e_{B_2} = e_{B_1 \uplus B_2}$ .

2.3. Flacets. Our last step will be to eliminate the redundancy from Theorem 1.10, i.e., to understand which flats do indeed define *facets* of  $P_M$  that meet the interior of  $r \cdot \Delta$ . For this, given a matroid M and a subset X of the ground set of M, aside from the already mentioned notion of restriction M[X], we need to define the "contraction" of X in M.

First recall that, if rk(X) = s, the bases of M[X] are

$$\mathcal{B}(M[X]) = \max \mathcal{I}[X] = \{X \cap B \mid B \in \mathcal{B}(M), |X \cap B| = s\}.$$

Now consider the set

$$\mathcal{B}_{/X} := \{ B \setminus X \mid B \in \mathcal{B}(M), |X \cap B| = s \}.$$

We would like to prove that  $\mathcal{B}_{/X}$  is the set of bases of a matroid. This we will do in Corollary 2.11, but first a word about duality.

**Lemma-Definition 2.10.** Let M be a matroid on the ground set E. Then

$$\mathcal{B}^*(M) := \{ E \setminus B \mid B \in \mathcal{B}(M) \}$$

is the set of bases of a matroid on E that is called dual to M, and denoted  $M^*$ .

*Proof.* Consider the following preparatory claim:

Let M be a matroid on the ground set E. Given  $B_1, B_2 \in \mathcal{B}(M)$  and  $e \in B_2 \setminus B_1$ , then there is  $f \in B_1 \setminus B_2$  such that  $(B_1 \setminus \{f\}) \cup \{e\} \in \mathcal{B}(M)$ .

*Proof.* By maximality of  $B_1$ , the set  $B_1 \cup \{e\}$  is dependent and thus contains a circuit C. Necessarily  $e \in C$ , otherwise  $C \subseteq B_1$  and the latter would be dependent. Also,  $C \neq \{e\}$ , since otherwise  $B_2$  would be dependent. Therefore  $C \cap B_1 \neq \emptyset$ , so we can choose f in this set and let  $B'_1 := B_1 \setminus \{f\} \cup \{e\}$ . Now  $B'_1$  is independent: any circuit  $C' \subseteq B'_1$  would contain e and, with (C3) applied to C, C' and e, it would lead to the existence of a circuit fully contained in  $B_1$ . Since  $|B_1| = |B'_1|$ , the claim follows.

Now consider  $B_1^*, B_2^* \in \mathcal{B}^*(M)$  and  $e \in B_1^* \setminus B_2^*$ . By definition,  $B_1 = E \setminus B_1^*$  and  $B_2 = E \setminus B_2^*$ are bases of M, and we have  $e \in B_2 \setminus B_1$ . Now by our "preparatory claim" we can find  $f \in B_1 \setminus B_2$ so that  $B_3 := (B_1 \setminus \{f\}) \cup \{e\}$  is in  $\mathcal{B}(M)$ . The basis  $B_3^* := E \setminus B_3 \in \mathcal{B}^*(M)$  then satisfies  $B_3^* = (B_1^* \setminus \{e\}) \cup \{f\}$  and witnesses ( $\mathcal{B}2$ ) for the chosen  $B_1^*, B_2^*, e$ . Aziom ( $\mathcal{B}1$ ) is obviously true for  $\mathcal{B}^*(M)$ , since this set contains the complement of any element of the (nonempty) set  $\mathcal{B}(M)$ . This concludes the proof.

M/X (Contraction)

**Corollary-Definition 2.11.** The set  $\mathcal{B}_{/X}$  is the set of bases of a matroid that we call M/X, the contraction of X in M.

*Proof.* This follows directly from Lemma 2.10 because  $\mathcal{B}_{/X} = \mathcal{B}^*(M^*[E \setminus X])$ .

To see this last fact, notice that the bases of  $M^*$  are the sets  $E \setminus B$  for B a basis of M; hence the bases of  $M^*[E \setminus X]$  are the sets  $(E \setminus B) \cap (E \setminus X)$  with maximal possible cardinality, i.e., such

(Dual matroid)

 $M^*$ 

that  $B \cap X$  has maximal cardinality, this maximal cardinality being  $s = \operatorname{rk}(X)$ . Therefore, writing  $X^c$  and  $B^c$  for  $E \setminus X$  and  $E \setminus B$ ,

$$\mathcal{B}^*(M^*[E \setminus X]) = \{\underbrace{X^c \setminus (B^c \cap X^c)}_{=X^c \setminus B^c = B \setminus X} \mid |B \cap X| = s\}$$

and the claim follows.

Flacet

**Definition 2.12.** A flat F of a matroid M is called a *flacet* if the inequality  $\sum_{i \in F} x_i \leq \operatorname{rk}(F)$  defines a facet of  $P_M$  that meets the interior of  $r \cdot \Delta$ .

**Proposition 2.13.** Let M be a connected matroid. A flat F is a flacet of M if and only if both M[F] and M/F are connected.

**Example 2.14.** See Figure 1 for an illustration of this theorem. All three "green" facets are flacets.

Proof of Proposition 2.13. Let F be a flat. The vertices of the face Q of  $P_M$  given by  $P_M \cap \{\sum_{i \in F} x_i = \operatorname{rk}(F)\}$  are exactly those  $e_B$  where  $B \in \mathcal{B}(M)$  is such that  $|B \cap F| = \operatorname{rk}(F)$ . Since F is a flat, these B are exactly (<sup>3</sup>) the bases of the matroid

$$M_F := M/F \oplus M[F].$$

Now, with Proposition 2.9 we can write

$$Q = P_{M_F} = P_{M/F} \times P_{M[F]}$$

and the dimension of Q can be computed as dim  $Q = \dim(P_{M/F}) + \dim(P_{M[F]}) = |E \setminus F| - c(M/F) + |F| - c(M[F]).$ 

Now Q is a facet if and only if dim  $Q = \dim P_M - 1 = n - 2$  (recall that M is connected). This is the case exactly when c(M/F) = c(M[F]) = 1.

<sup>&</sup>lt;sup>3</sup>The inclusion  $\{B \in \mathcal{B}(M) \mid |B \cap F| = \operatorname{rk}(F)\} \subseteq \mathcal{B}(M/F \oplus M[F])$  is clear. For the reverse inclusion consider any basis  $B_1 \oplus B_2$  of  $M/F \oplus M[F]$  and notice that if it not a basis of M then there must be some  $e \in B_1 \subseteq E \setminus F$  such that  $\{e\} \cup B_2$  is dependent, i.e.,  $\operatorname{rk}(\{e\} \cup B_2) = |B_2| = \operatorname{rk}(F)$ . Now (R3) gives  $\operatorname{rk}(B_2 \cup \{e\}) + \operatorname{rk}(F) \geq \operatorname{rk}(B_2) + \operatorname{rk}(F \cup \{e\})$  and so we conclude  $\operatorname{rk}(F) \geq \operatorname{rk}(F \cup \{e\})$ , which implies  $\operatorname{rk}(F) = \operatorname{rk}(F \cup \{e\})$  via (R2), a contradiction to F being a flat.